

#### Yury Lifshits<sup>1,2</sup> and Dirk Nowotka<sup>3</sup>

<sup>1</sup>Steklov Institute of Mathematics at St.Petersburg, <sup>2</sup>California Institute of Technology, <sup>3</sup>University of Stuttgart



#### Y. Lifshits, D. Nowotka

Estimation of Click Volume

CSR 2007 1 / 20

## Part I

## Fast Algorithm for Solving Least Squares on Sparse Matrices

#### General area:

Fast web-scale algorithms for machine learning

Least Squares: Find  $\alpha$  minimizing  $\|\alpha M - y\|_2$ 

Assume *M* has  $k \ll mn$  nonzero elements Can we solve least squares in O(k) time?

#### **Our contribution:**

- Solving least squares on sparse matrices
- Application to on-line advertisement

## Geometric View on Least Squares

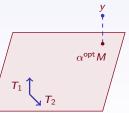
Find  $\alpha$  minimizing  $\|\alpha M - y\|_2$ 

Let V be a linear hull of M's rows  $T_1, \ldots, T_m$ 

Interpretation: every vector  $\alpha$  represents coordinates of  $\alpha M$  point in V in  $T_1, \ldots, T_m$  basis

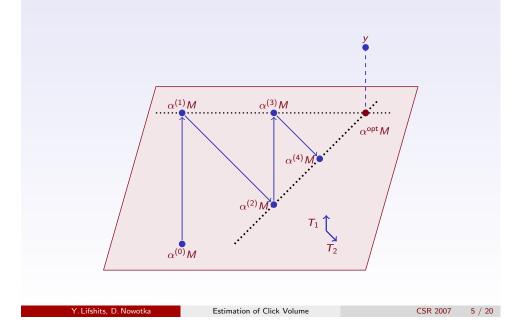
Estimation of Click Volu

Thus  $\alpha^{\text{opt}} M$  is just the projection of y to V

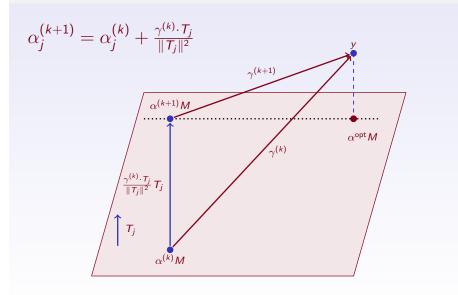


CSR 2007

## Componentwise Iterations: Idea



## Optimality of One Step Update



## Componentwise Iterations: Algorithm

#### Auxiliary data structures:

Matrix columns:  $T_1, \ldots, T_m$ , precomputed norms  $||T_j||$ Current solution  $\alpha^{(k)}$ Error vector for current prediction  $\gamma^{(k)} = y - \alpha^{(k)}M$ 

#### **Componentwise iterations:**

- Start with random  $\alpha^{(0)}$
- **2** Choose some coordinate j. Update formulae:

Estimation of Click Volum

 $\alpha_j^{(k+1)} = \alpha_j^{(k)} + \frac{\gamma^{(k)} \cdot T_j}{\|T_j\|^2}$  $\gamma^{(k+1)} = \gamma^{(k)} - \frac{\gamma^{(k)} \cdot T_j}{\|T_j\|^2} T_j$ 

CSR 2007 6 / 20

## Convergence Theorem

Theorem (Convergence theorem)  $\alpha^{(k)}$  converges to  $\alpha^{opt} = (M \cdot M^T)^{-1}M^T y$  for any infinite order of updates containing every j infinitely many times

**Open Problem:** prove some upper bounds on convergence speed

## Complexity of One Global Round

#### **Global Round:**

Sequentially do one update for every j from 1 to m

#### Theorem

Componentwise Iterations algorithm uses only O(k) time for global round of updates. Recall, k is the number of nonzero entries in M.



## Algorithm: Discussion

- A vector α<sup>(k)</sup> can be safely updated in two components j<sub>1</sub> and j<sub>2</sub> in parallel if we have T<sub>j1</sub> ⊥ T<sub>j2</sub>
- In the case of orthogonal vectors  $T_1, \ldots, T_m$  one global round is sufficient for reaching  $\alpha^{\text{opt}}$
- Joint optimal update of k components requires inverting  $k \times k$  matrix

Estimation of Click Volu

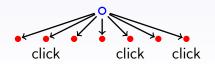
## On-line Ads: Terminology

Opportunity: all available information about ad request to advertising system. E.g. website, search query, user id, . . .

Click-through rate CTR(a, o): a probability of clicking ad a at opportunity o

Click volume for collection *O* of all opportunities for some time interval:  $CV(a) = \sum_{o \in O} CTR(a, o)$ 

The same ad to all ad requests



CSR 2007

## Importance of Click Volume

- Measuring effectiveness of advertising system by a ration between sold clicks and click volume
- Measuring preciseness of targeting by ration between click volume in target group and overall click volume
- Defining coverage/agressiveness tradeoff: How many users should we contact in order to get a given fraction of click volume?

Estimation of Click Volue

• Helpful for setting prices for on-line ads

### Estimation of Click Volume: Methodology

- Indexing: mapping advertising system logs to well-defined data structure
- Regression Analysis: deriving formula for click-through rate prediction
- Estimation: applying resulting formula to a given ad and monthly collection of opportunities

Estimation of Click Volu

## Indexing: History Table

#### Logs of advertising system:

Set of triples  $(a_i, o_i, b_i)$ : ad, opportunity, click

#### **Events indexing:**

Mapping a pair (a, o) to internal representation E(a, o)Techniques: clustering, dimensionality reduction

#### **Resulting reduced table:**

Set of pairs  $(E_i, CTR_i)$ 

## Quick Recall: Regression Analysis

#### Input:

Training collection of n documents Document i: m-dimensional vector  $D_i$ and additional parameter  $y_i$ 

#### **Regression Problem:** Find function f such that $f(D_i) \approx y_i$

**Linear Regression Problem** (least squares): find *m*-dimensional vector  $\alpha$  such that the sum of squared prediction errors  $\sum |\alpha D_i - y_i|^2$  is minimized

CSR 2007

CSR 2007

## Click Volume via Regression

- Building history table from logs of advertising system: matrix  $M = \{E_1, \dots, E_n\}$  and *CTR* vector
- Solving least squares
- Omputing click volume by formula:

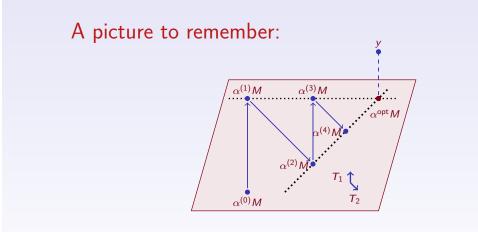
## $CV(a_{\text{new}}) = \sum_{1 \le i \le n} \alpha \cdot E(a_{\text{new}}, o_i)$

Y. Lifshits, D. Nowotka

Estimation of Click Volume

CSR 2007

17 / 20



Thanks for your attention! Questions?

## Directions for Further Work

#### **Theoretic problems:**

- Prove upper bounds for speed of convergence of our algorithm
- Can we compute click volume for all ads in the system faster than doing it separately for every ad?

#### **Experimental problems:**

- Apply our algorithm for some real data set. Measure the empirical precision of *CTR* prediction
- Study effects of heuristical ingredients for algorithm: indexing, dimensionality reduction, update order

Estimation of Click Volume

# Yury Lifshitshttp://yury.nameDirk Nowotkagoogle://dirk nowotka

#### Some related work:

/ Lifshits D Nowotka

- Y. Lifshits, D. Nowotka Estimation of the Click Volume by Large Scale Regression Analysis. CSR'07. http://yury.name/papers/lifshits2007click.pdf
- B. Hoffmann, Y. Lifshits, D. Nowotka Maximal Intersection Queries in Randomized Graph Models. CSR'07. http://yury.name/papers/hoffmann2007maximal.pdf
- N. Goyal, Y. Lifshits, H. Schütze Disorder Inequality: A Combinatorial Approach to Nearest Neighbor Search. Submitted. http://yury.name/papers/goyal2008disorder.pdf

CSR 2007