# Maximal Intersection Queries in Randomized Graph Models 

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> CSBI

## The Maximal Intersection Problem

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Database: A family $\mathcal{F}$ of $n$ sets, $|f| \leq k \forall f \in \mathcal{F}$.
Query: Given a set $f_{\text {new }}$ with $\left|f_{\text {new }}\right| \leq k$, return $f_{i} \in \mathcal{F}$ with maximal $\left|f_{\text {new }} \cap f_{i}\right|$.
Constraints: Preprocessing time $n \cdot p o l y l o g(n) \cdot p o l y(k)$. Query time polylog(n) • poly (k).


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## Motivation

## Possible applications of MaxInt:

- Advertisement matching: Which website should an advertisement be placed on (e.g. Google AdSense)?
- Text clustering/classification: Find a document in a database that has a maximal number of common terms with a newcomer document (e.g. Reuters).
- Recommendation systems, near-duplicate detection, code plagiarism detection, search engines, ...


## Relation to Nearest Neighbor

Nearest Neighbor Problem: Determine in a general metric space a point that is closest to a given query point (Zezula et al., Similarity Search - The Metric Space Approach, Springer, 2006).

MaxInt: special case of Nearest Neighbor.
Similarity: size of intersection.

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## Main result

Assumption: Input is taken from some predefined distribution.

There exists an algorithm that finds with very high probability an almost optimal solution in time logarithmic in the size of the family.

Zipf ${ }^{1}$ (1932): In natural language texts the (absolute) frequency $f$ of a term is approximately inversely proportional to its rank $r$ in the frequency table.
$\exists$ constant $c$ such that $f \cdot r \approx c$.
${ }^{1}$ George Kingsley Zipf, 1902-1950

## Zipf's law

| Word | Freq. | Rank | $f \cdot r$ | Word | Freq. | Rank | $f \cdot r$ |
| :--- | ---: | ---: | ---: | :--- | ---: | ---: | ---: |
| the | 3332 | 1 | 3332 | turned | 51 | 200 | 10200 |
| and | 2972 | 2 | 5944 | you'll | 30 | 300 | 9000 |
| a | 1775 | 3 | 5235 | name | 21 | 400 | 8400 |
| he | 877 | 10 | 8770 | comes | 16 | 500 | 8000 |
| but | 410 | 20 | 8400 | group | 13 | 600 | 7800 |
| be | 294 | 30 | 8820 | lead | 11 | 700 | 7700 |
| there | 222 | 40 | 8880 | friends | 10 | 800 | 8000 |
| one | 172 | 50 | 8600 | begin | 9 | 900 | 8100 |
| about | 158 | 60 | 9480 | family | 8 | 1000 | 8000 |
| more | 138 | 70 | 9660 | brushed | 4 | 2000 | 8000 |
| never | 124 | 80 | 9920 | sins | 2 | 3000 | 6000 |
| Oh | 116 | 90 | 10440 | Could | 2 | 4000 | 8000 |
| two | 104 | 100 | 10400 | Applausive | 1 | 8000 | 8000 |

Empirical evaluation of Zipf's law on Tom Sawyer (Manning/Schütze, Foundations of statistical natural language processing, MIT Press, 1999)

## Zipf model

Documents $\mathcal{D}=\left\{d_{1}, \ldots, d_{n}\right\}$ Terms $\mathcal{T}=\left\{t_{1}, \ldots, t_{m}\right\}, m \leq \operatorname{poly}(n)$

## Generating a document collection:

- Every document is generated independently.
- Term occurences are also independent.
- A document contains term $t_{i}$ with probability $\frac{1}{i}$.
- Expected number of terms in a document: $\ln m$.

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## Zipf model (cont.)

Definition (Relative frequency of a term $t$ in a document collection D):

$$
\frac{|\{d \in \mathcal{D} \mid t \in d\}|}{|\mathcal{D}|}
$$

In the Zipf model:

- Expected relative frequency for $t_{i}$ : $\frac{1}{i}$
- Expected frequency rank for $t_{i}$ : $i$-th value among those of all terms.
$\Rightarrow$ Zipf model reflects in a natural way Zipf's law.


## Zipf model: relation to random graphs

Representing the data as a graph: bipartite graph (documents/terms)

$\Rightarrow$ Zipf's law for distribution of term degrees, $P(k)=\frac{1}{k}$.

## Threshold phenomenon

Assume that the terms of a query document are ordered by their frequency in the document collection.

Two events:

- Any $q$-match: $\exists d \in D$ that has at least $q$ common terms with the query document.
- Prefix $q$-match: $\exists d \in D$ that has at least $q$ "top" terms with the query document.

For a document collection following the Zipf model the following holds:

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## Threshold phenomenon



Exemplary probability curves for any $q$-match and prefix $q$-match. ( $q=$ no. of matched terms)

## Zipf model (cont.)

Partitioning of $\mathcal{T}$ :


Group $P_{i}$ includes terms from $t_{\left[e^{i-1}\right\rceil}$ to $t_{\left\lfloor e^{i}\right\rfloor}$.
A document that contains $\ln m$ terms $p_{1} \ldots p_{\text {ln } m}, p_{i} \in P_{i}$, will be called regular.

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Partitioning of $\mathcal{T}$ :


Group $P_{i}$ includes terms from $t_{\left\lceil e^{i-1}\right\rceil}$ to $t_{\left\lfloor e^{i}\right\rfloor}$.
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## Magic level:

$$
q=\sqrt{2 \ln n} \quad(n=\text { no. of documents }) .
$$

## Theorem

Let $3 \leq \gamma<q-1, \gamma \in \mathbb{N}$. Fix $n, m$ and a regular query document $d_{\text {new }}$. Then for a document collection following the Zipf model the following holds:
(1) The probability that there exists a document $d \in \mathcal{D}$ that contains the first $q-\gamma$ terms of $d_{\text {new }}$ is greater than $1-2^{-\frac{q(\gamma+1)}{2}}$.
(2) The probability that there exists a document $d \in \mathcal{D}$ that contains at least $q+\gamma$ terms of $d_{\text {new }}$ is smaller than $\frac{1}{e^{(\gamma-2) q-1}}$.

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## MAXINT algorithm in the Zipf model

## Preprocessing

(1) For every document: Sort the term list according to the position of the term in the frequency table.
(2) For every document: Generate the set of all possible regular $(q-\gamma)$-lists.
(3) Sort these regular lists and store for every list a pointer to the corresponding document.
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Complexity: $\mathcal{O}(\log m \cdot n \cdot \log n)$
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