# Open Problems TO GO 

Yury Lifshits<br>Caltech<br>http://yury.name



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- Short mathematical statement
- No background required
- Motivation (importance) is guaranteed

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Today:
Three classic problems
Three problems from YL

## 1 <br> Classic Problems

Mean payoff games
Semi-Thue systems
Ulam conjecture (graph reconstruction)

# 1.1. Rules of mean payoff games 

Input for a mean payoff game:

- Weighted directed graph (integer weights)
- Graph does not contain simple cycles with zero sum
- Vertices are divided into disjoint sets $A$ and $B$
- The starting vertex


## Rules of Mean Payoff Games

- Two players: Alice and Bob
- Players move the token over arcs
- Game starts from the starting vertex and it is infinite
- Alice plays from vertices of $A$, Bob from these of $B$
- Alice wins if the sum of already passed arcs goes to +infty
- Bob wins if the sum of already passed arcs goes to -infty


## Computational Problem

Given a game graph with an $A, B$ decomposition and a starting vertex to determine the winner (and find the winning strategy)

# MPG is Very Challenging 

MPG Problem belongs to NPnco-NP Direct applications in $\mu$-calculus verification

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Known algorithms:

- Randomized algorithm $\mathcal{O}^{*}\left(2^{\sqrt{n}}\right)$ expected time
- Deterministic algorithm $\mathcal{O}^{*}\left(2^{n}\right)$ time


## References

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Journal of Mathematical Sciences, 2007
http://yury.name/papers/lifshits2006fast.pdf
M. Jurdziński, M. Paterson, U. Zwick

A deterministic subexponential algorithm for solving parity games
SODA'06
http://www.dcs.warwick.ac.uk/ ${ }^{\sim} m j u / P a p e r s / J P Z 07$-manuscript.pdf

H. Björklund, S. Vorobyov

A combinatorial strongly subexponential strategy improvement algorithm for mean payoff games

Discrete Applied Mathematics, 2007
http://portal.acm.org/citation.cfm?id=1222484

## Ulam Conjecture

A vertex-deleted subgraph of a graph $G$ is a subgraph $G-v$ obtained by deleting a vertex $v$ and its incident edges. The deck of a graph $G$ is the family of (unlabelled) vertex-deleted subgraphs of $G$; these are the cards of the deck. A reconstruction of a graph $G$ is a graph $H$ with the same deck as $G$. A graph $G$ is reconstructible if every reconstruction of $G$ is isomorphic to $G$.

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Conjecture: every graph with at least three vertices is reconstructible

## Reference

J.A. Bondy

A graph reconstructor's manual
Surveys in Combinatorics, 1991
http://www.ecp6.jussieu.fr/pageperso/bondy/research/papers/recon.ps

## Semi-Thue Systems

Rewriting $(\alpha, \beta)$ rule allows to rewrite any $u \alpha v$ in $u \beta v$

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Word problem: Given system of rules and two words $w_{1}$ and $w_{2}$ to decide whether one can be obtained from another by a sequence of such rules?

## Challenge

There is a system with three rules such that word problem is undecidable

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Is word problem decidable or not for systems of one (two) rules?

## Reference

Y. Matiyasevich and G. Senizerguez

Decision Problems for Semi-Thue Systems with a Few Rules
LICS'96
http://dept-info.labri.u-bordeaux.fr/~ges/termination.ps

## Open Problems from YL

## Compressed Arithmetic

Input: Two grammars of size $n, m$ generating binary strings $P$ and $Q$ of the same length

Task: Compute a close-to-minimal grammar generating "bitwise OR between $P$ and $Q$ "

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Can we do it in time poly $(n+m+$ output $)$ ?

## References

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Processing Compressed Texts: A Tractability Border
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http://yury.name/papers/lifshits2007processing.pdf


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Querying and Embedding Compressed Texts
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Window Subsequence Problems for Compressed Texts
CSR'06
http://yury.name/papers/cegielski2006window.pdf

# Impossibility of Preprocessing 

## Input

Circuits $C_{1} \ldots, C_{n}$
of size poly $(m)$ with input size $m$
Query task
Given string $y$ of length $m$ to answer whether $\exists i: C_{i}(y)=y e s$

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Open problem: Is there a solution within given constraints?

## Dual Problem

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Given circuit $C$ of size polym with input length $m$ to answer whether $\exists i: C\left(x_{i}\right)=y e s$

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## Reference

Yury Lifshits
Algorithms for Nearest Neighbors: Classic Ideas, New Ideas
Talk at University of Toronto
MP3 recording
http://yury.name/talks/toronto-talk.pdf

## Positive Subgraph

## Input <br> $n \times n$ bipartite graph (pretty sparse) Weights on edges

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Polynomial approximate algorithm?

## Reference

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Y. Lifshits and D. Nowotka

Estimation of the click volume by large scale regression analysis
CSR'07
http://yury.name/papers/lifshits2007click.pdf
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## Voting

## Which problem you like the most?

- Mean Payoff Games
- Ulam Conjecture
- Semi-Thue Systems
- Compressed Arithmetics
- Impossibility of Preprocessing
- Positive Subgraph
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## Thanks for your attention! Questions?

